



The wonderful life of a SQL query in a streaming database



RisingWave Labs

- Creator of the RisingWave Database
- OLAP streaming queries
- Incremental updates on materialized views



Views and Materialized Views

- What is a view?
- What is a MV?
- How do you calc views traditionally?
 - Full rebuild is expensive
- Incremental updates
 - Run aggregate on diff
 - Run background job that detects changes in base table
 - Triggers that fire if there is a change to the underlying table
- RisingWave: Incremental updates

```
... ⌘1 CREATE_TABLE
CREATE_MATERIALIZED_VIEW
dev=> INSERT INTO stories (id, author, title, url) VALUES (1, 2, 'hacker story', 'some-url.net');
INSERT 0 1
dev=> SELECT * FROM stories;
 id | author | title | url
-----+-----+-----+-----
  1 |      2 | hacker story | some-url.net
(1 row)

dev=> INSERT INTO votes (user, story_id) VALUES (2, 1), (3, 1);
INSERT 0 2
dev=> SELECT * FROM votes;
 user | story_id
-----+-----
    3 |         1
    2 |         1
(2 rows)

dev=> SELECT * FROM StoriesWithVC;
 id | author | title | url | vcount
-----+-----+-----+-----+-----
  1 |      2 | hacker story | some-url.net | 2
(1 row)


dev=>
dev=>
```



Streaming graph


```
CREATE TABLE stories (id int, author int, title text, url text);  
CREATE TABLE votes (user int, story_id int);
```

```
CREATE MATERIALIZED VIEW StoriesWithVC AS  
SELECT id, author, title, url, vcount  
FROM stories  
JOIN ( SELECT story_id, COUNT(*) AS vcount FROM votes GROUP BY story_id) as VoteCount  
on VoteCount.story_id = stories.id;
```



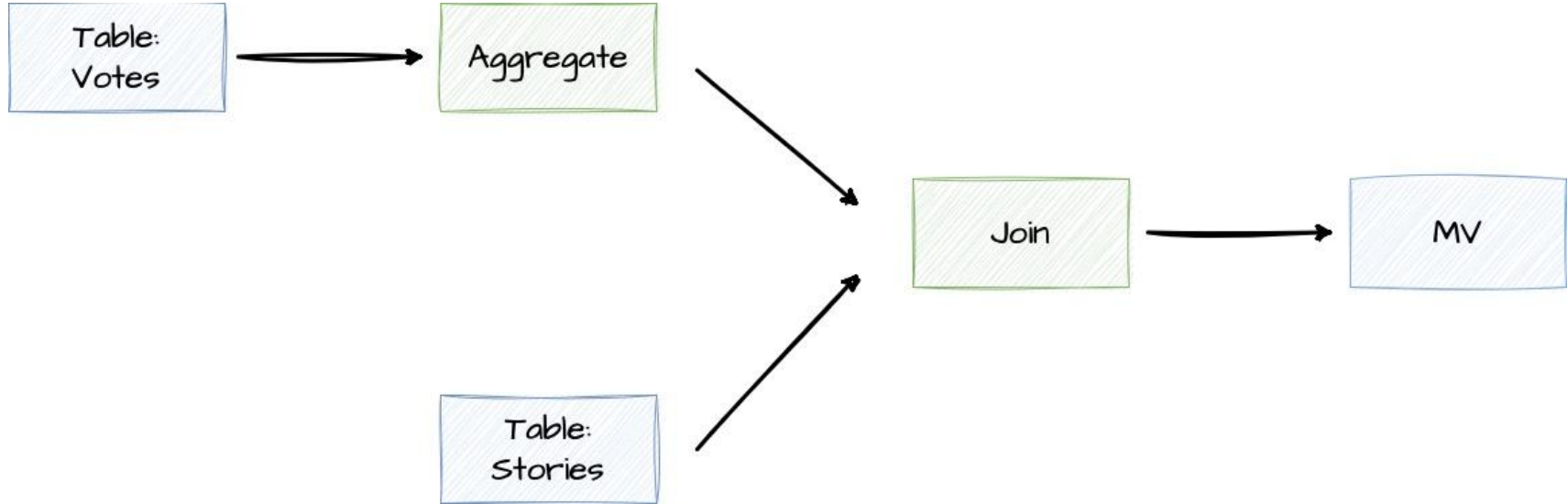
```
EXPLAIN CREATE MATERIALIZED VIEW StoriesWithVC AS
SELECT id, author, title, url, vcount
FROM stories JOIN (SELECT story_id, COUNT(*) AS vcount FROM votes GROUP BY story_id) AS VoteCount
ON VoteCount.story_id = stories.id;
```

```
StreamMaterialize { columns: [id, author, title, url, vcount, ...] }
└─StreamExchange
  └─StreamHashJoin { type: Inner, predicate: stories.id = votes.story_id }
    └─StreamExchange { dist: HashShard(stories.id) }
      └─StreamTableScan { table: stories, columns: [id, author, title, url, _row_id] }
        └─StreamHashAgg { group_key: [votes.story_id], aggs: [count] }
          └─StreamExchange { dist: HashShard(votes.story_id) }
            └─StreamTableScan { table: votes, columns: [story_id, _row_id] }
```

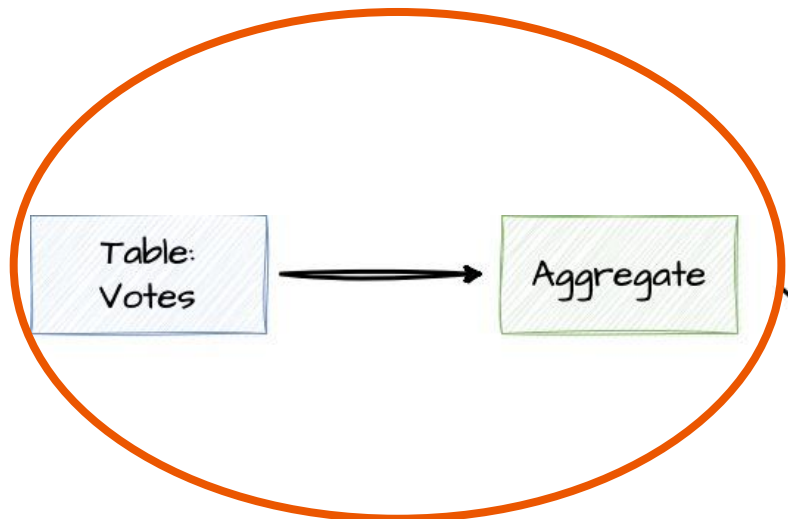


```
EXPLAIN CREATE MATERIALIZED VIEW StoriesWithVC AS
SELECT id, author, title, url, vcount
FROM stories JOIN (SELECT story_id, COUNT(*) AS vcount FROM votes GROUP BY story_id) AS VoteCount
ON VoteCount.story_id = stories.id;
```

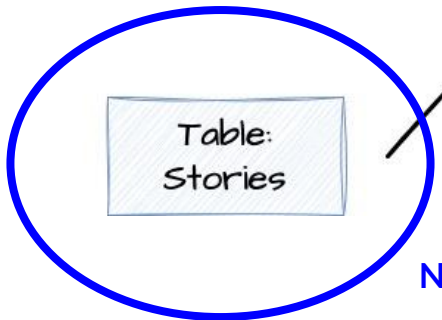
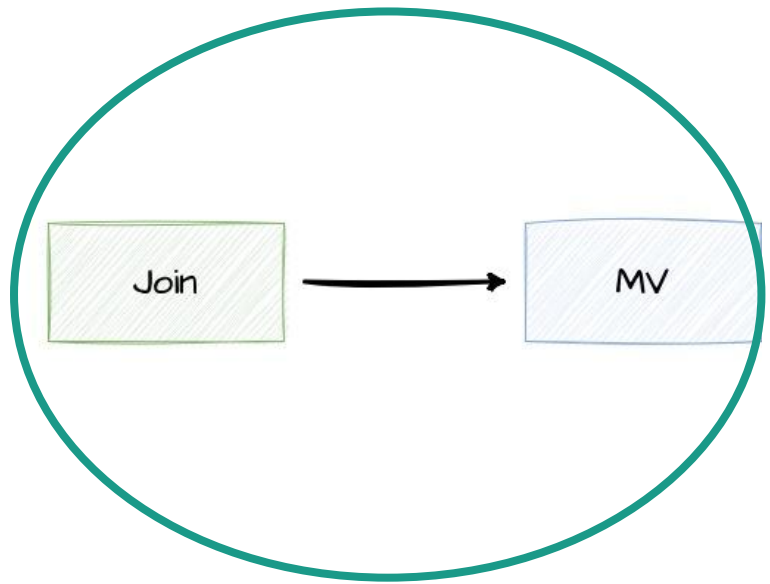
```
StreamMaterialize { columns: [id, author, title, url, vcount, ...] }
└─StreamExchange
  └─StreamHashJoin { type: Inner, predicate: stories.id = votes.story_id }
    └─StreamExchange { dist: HashShard(stories.id) }
      └─StreamTableScan { table: stories, columns: [id, author, title, url, _row_id] }
        └─StreamHashAgg { group_key: [votes.story_id], aggs: [count] }
          └─StreamExchange { dist: HashShard(votes.story_id) }
            └─StreamTableScan { table: votes, columns: [story_id, _row_id] }
```



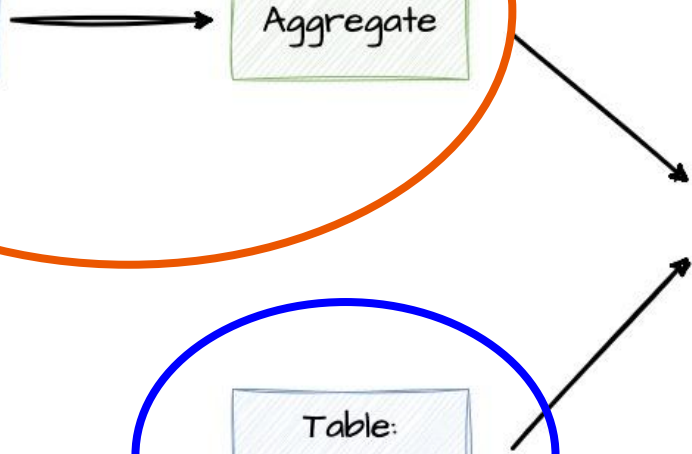
Node 1

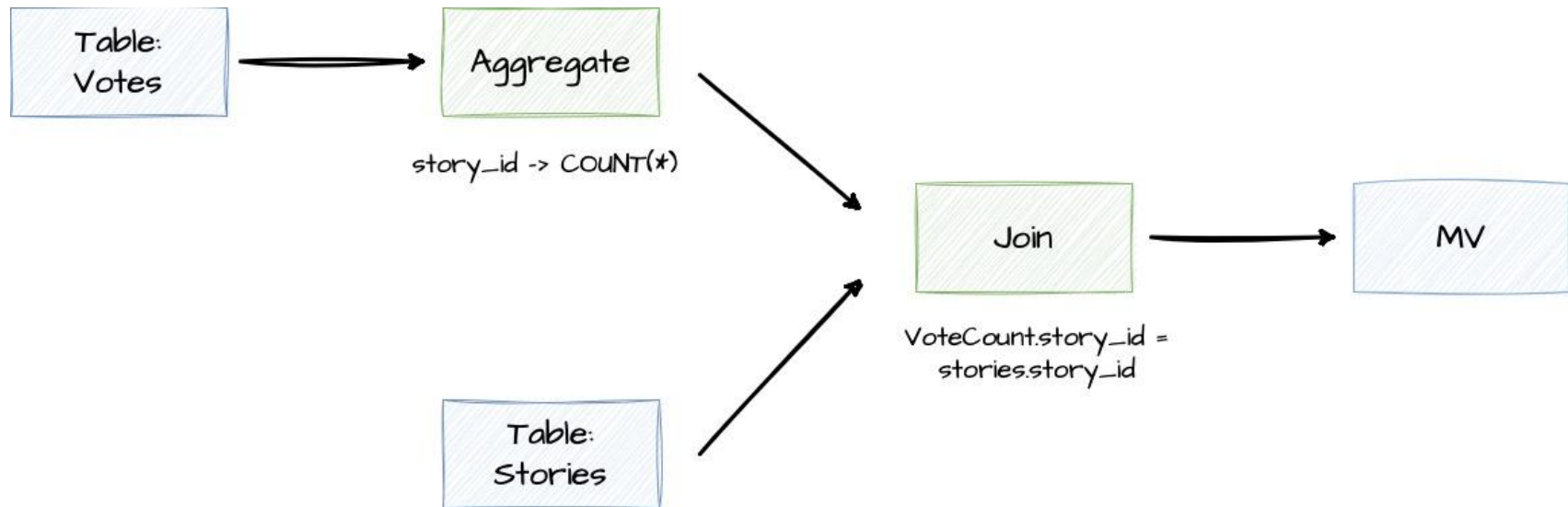


Node 2



Node 3





user	story_id
1234	1003
2345	1004
9999	1003

Table:
Votes

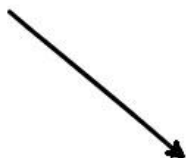


Aggregate

story_id -> COUNT(*)

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories



Join

VoteCount.story_id =
stories.story_id



MV

user	story_id
1234	1003
2345	1004
9999	1003

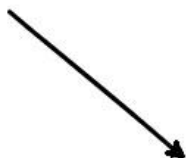
story_id	COUNT(*)
1003	2
1004	1

Table:
Votes



Aggregate

story_id -> COUNT(*)



Join



MV

VoteCount.story_id =
stories.story_id

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories



user	story_id
1234	1003
2345	1004
9999	1003

story_id	COUNT(*)
1003	2
1004	1

Table:
Votes

Aggregate

story_id -> COUNT(*)

story_id	COUNT(*)
1003	2
1004	1

ID	author
1003	Lu
1004	Jane
1005	John

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories

Join

VoteCount.story_id =
stories.story_id

MV

user	story_id
1234	1003
2345	1004
9999	1003

story_id	COUNT(*)
1003	2
1004	1

Table:
Votes

Aggregate

story_id -> COUNT(*)

story_id	COUNT(*)
1003	2
1004	1

ID	author
1003	Lu
1004	Jane
1005	John

Join

MV

VoteCount.story_id =
stories.story_id

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories

ID	author	title	url	vcount
1003	Lu	foo	http	2
1004	Jane	bar	http	1

user	story_id
1234	1003
2345	1004
9999	1003
789	1004

story_id	COUNT(*)
1003	2
1004	1

Table:
Votes

Aggregate

story_id -> COUNT(*)

story_id	COUNT(*)
1003	2
1004	1

ID	author
1003	Lu
1004	Jane
1005	John

Join

VoteCount.story_id =
stories.story_id

MV

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories

ID	author	title	url	vcount
1003	Lu	foo	http	2
1004	Jane	bar	http	1

user	story_id
1234	1003
2345	1004
9999	1003
789	1004

story_id	COUNT(*)
1003	2
1004	1

Table:
Votes

Aggregate

+ (789, 1004)

story_id	COUNT(*)
1003	2
1004	1

ID	author
1003	Lu
1004	Jane
1005	John

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories

Join

MV

ID	author	title	url	vcount
1003	Lu	foo	http	2
1004	Jane	bar	http	1

user	story_id
1234	1003
2345	1004
9999	1003
789	1004

story_id	COUNT(*)
1003	2
1004	2

Table:
Votes

Aggregate

story_id	COUNT(*)
1003	2
1004	1

ID	author
1003	Lu
1004	Jane
1005	John

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories

Join

MV

ID	author	title	url	vcount
1003	Lu	foo	http	2
1004	Jane	bar	http	1

user	story_id
1234	1003
2345	1004
9999	1003
789	1004

story_id	COUNT(*)
1003	2
1004	2

story_id	COUNT(*)
1003	2
1004	1

ID	author
1003	Lu
1004	Jane
1005	John

Table:
Votes

Aggregate

- (1004, 1)
+ (1004, 2)

Join

MV

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories

ID	author	title	url	vcount
1003	Lu	foo	http	2
1004	Jane	bar	http	1

user	story_id
1234	1003
2345	1004
9999	1003
789	1004

story_id	COUNT(*)
1003	2
1004	2

Table:
Votes

Aggregate

story_id	COUNT(*)
1003	2
1004	2

ID	author
1003	Lu
1004	Jane
1005	John

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories

Join

MV

ID	author	title	url	vcount
1003	Lu	foo	http	2
1004	Jane	bar	http	1

user	story_id
1234	1003
2345	1004
9999	1003
789	1004

story_id	COUNT(*)
1003	2
1004	2

Table:
Votes

Aggregate

story_id	COUNT(*)
1003	2
1004	2

ID	author
1003	Lu
1004	Jane
1005	John

Join

MV

- (1004, Jane, bar, http, 1)
+ (1004, Jane, bar, http, 2)

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories

ID	author	title	url	vcount
1003	Lu	foo	http	2
1004	Jane	bar	http	1

user	story_id
1234	1003
2345	1004
9999	1003
789	1004

story_id	COUNT(*)
1003	2
1004	2

Table:
Votes

Aggregate

story_id	COUNT(*)
1003	2
1004	2

ID	author
1003	Lu
1004	Jane
1005	John

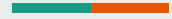
ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories

Join

MV

ID	author	title	url	vcount
1003	Lu	foo	http	2
1004	Jane	bar	http	2



Distributed systems



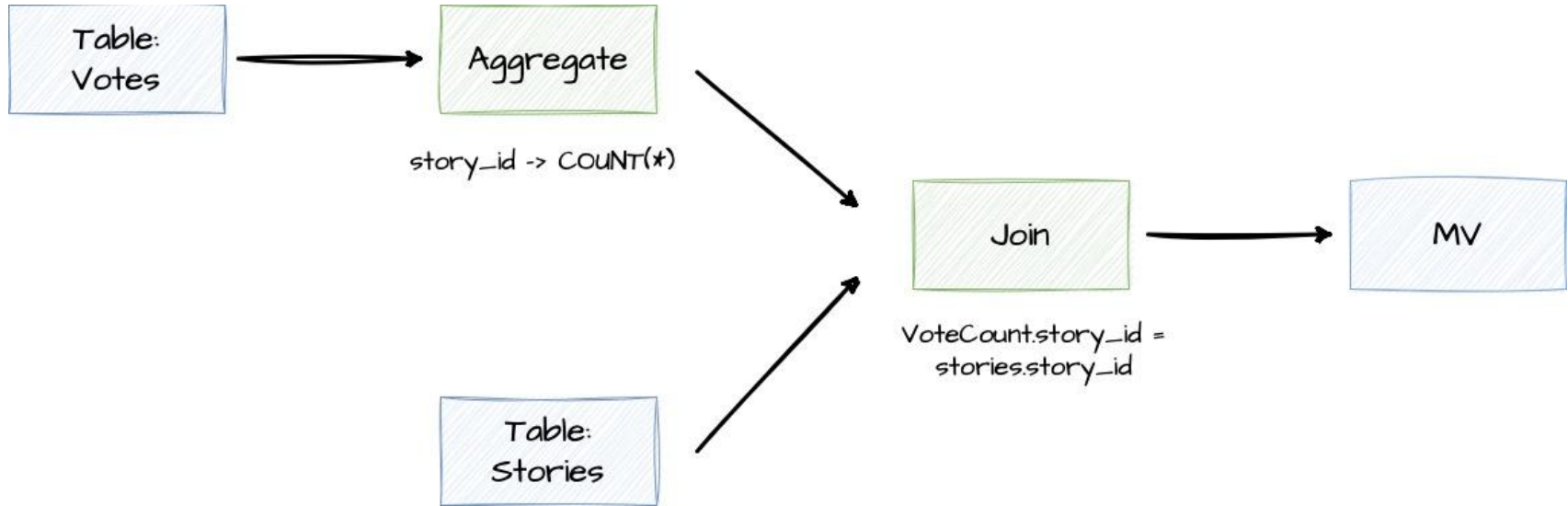
Challenges and opportunities in a distributed setup

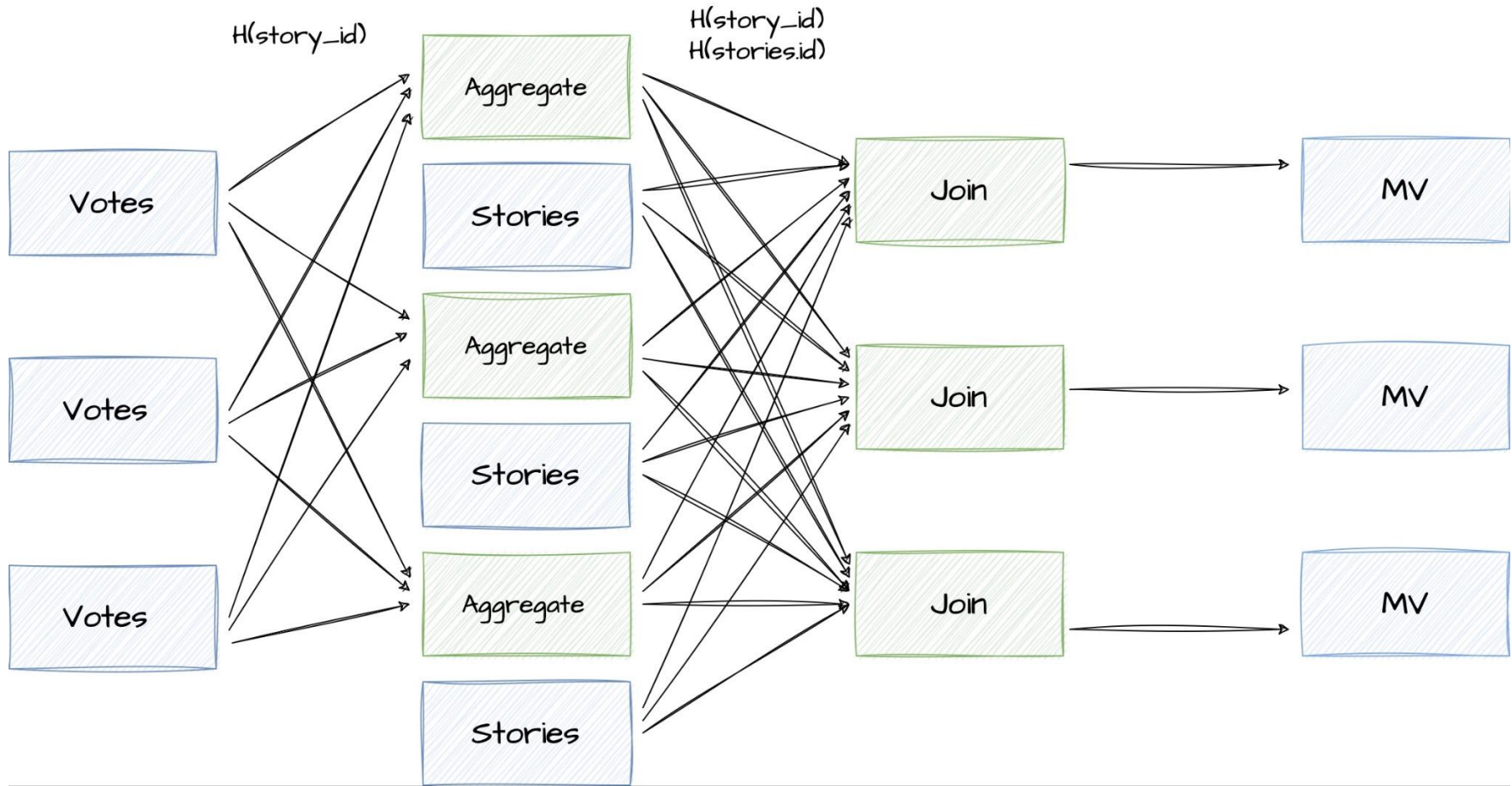
- Opportunity:
 - Execute in parallel
- Challenges:
 - **Recovery:** Trying not to lose data when a node crashes
 - **Scalability:** Adding/removing nodes if you have more/less workloads

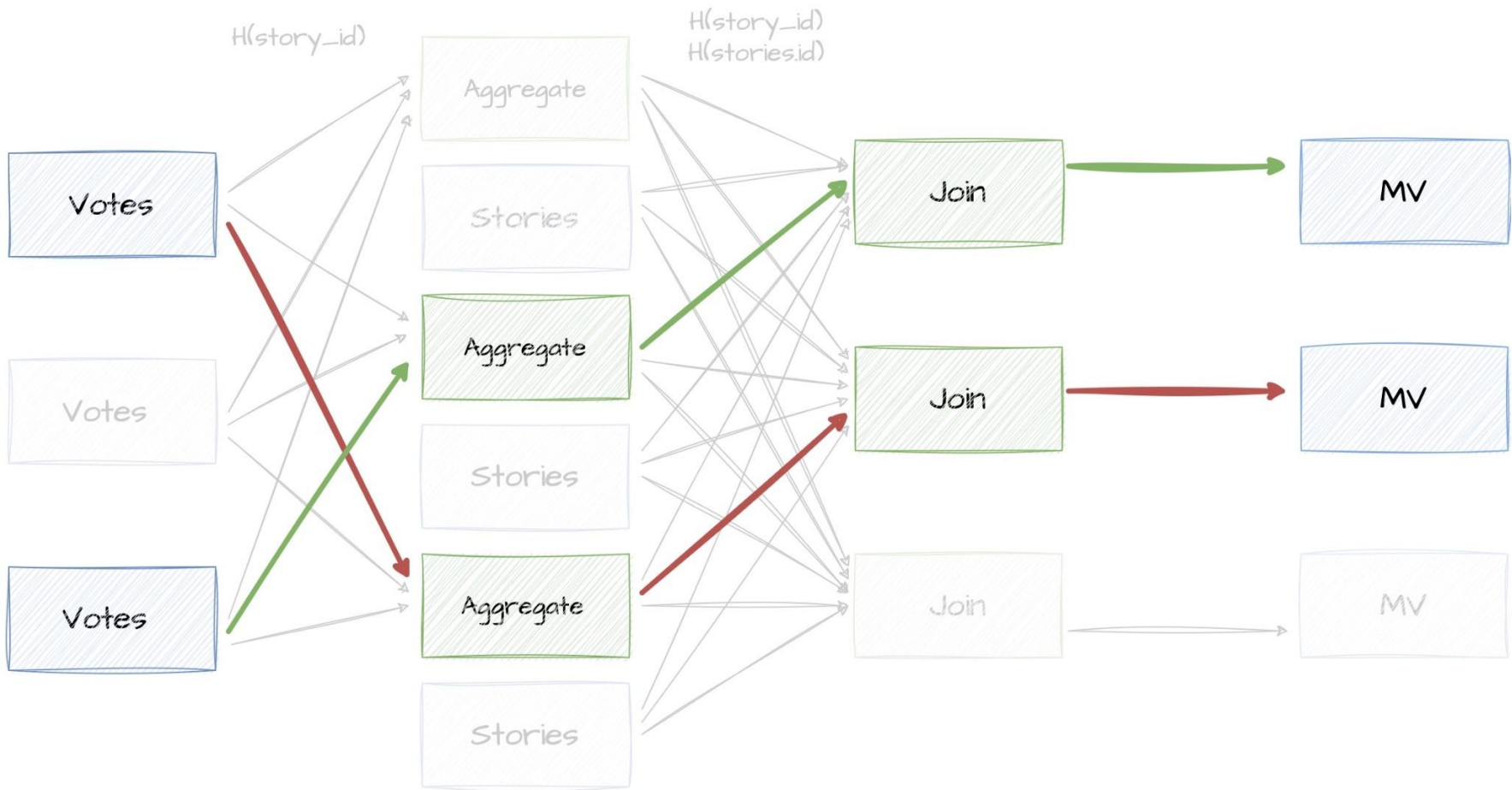


Distributed systems

Parallelism

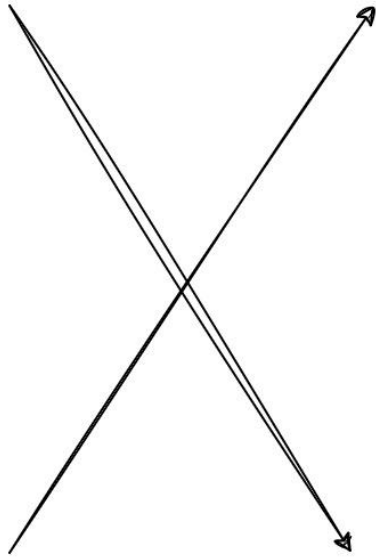






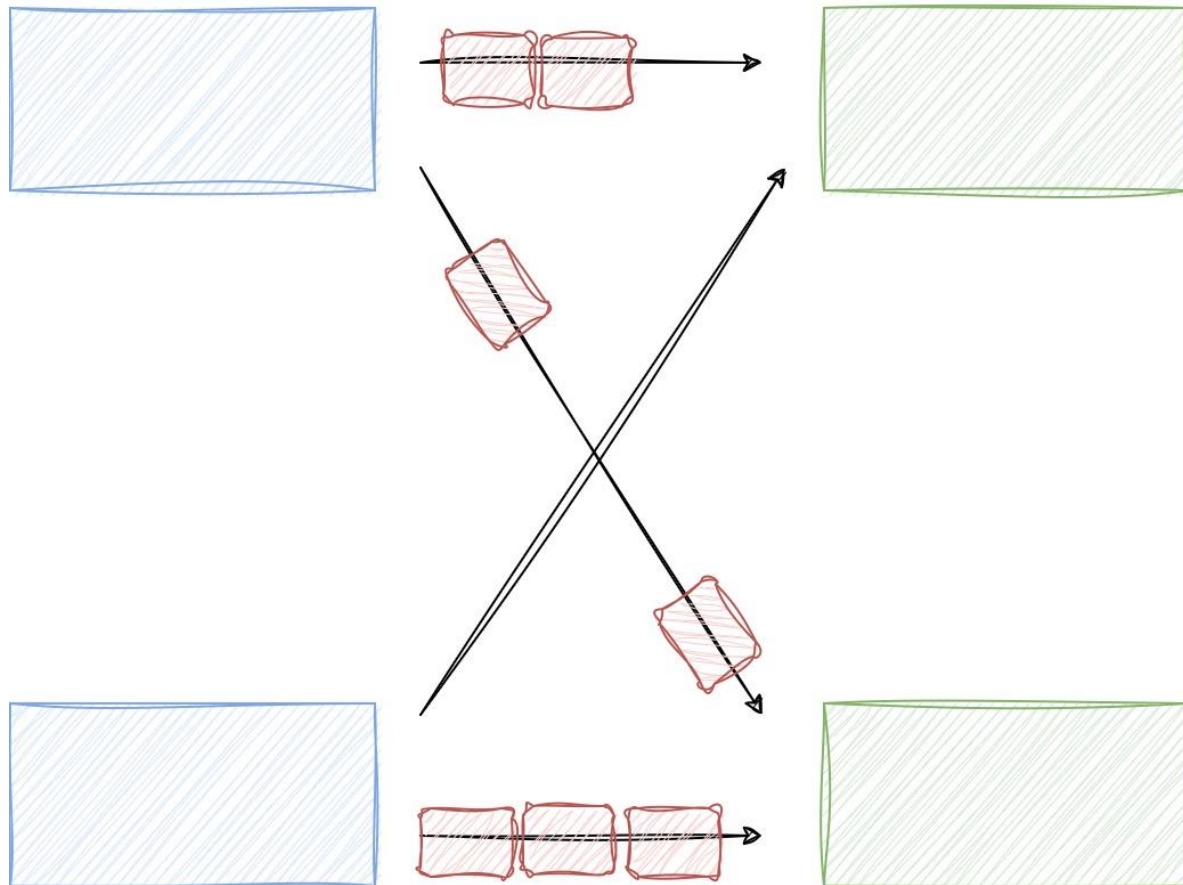
Scan

Aggregate



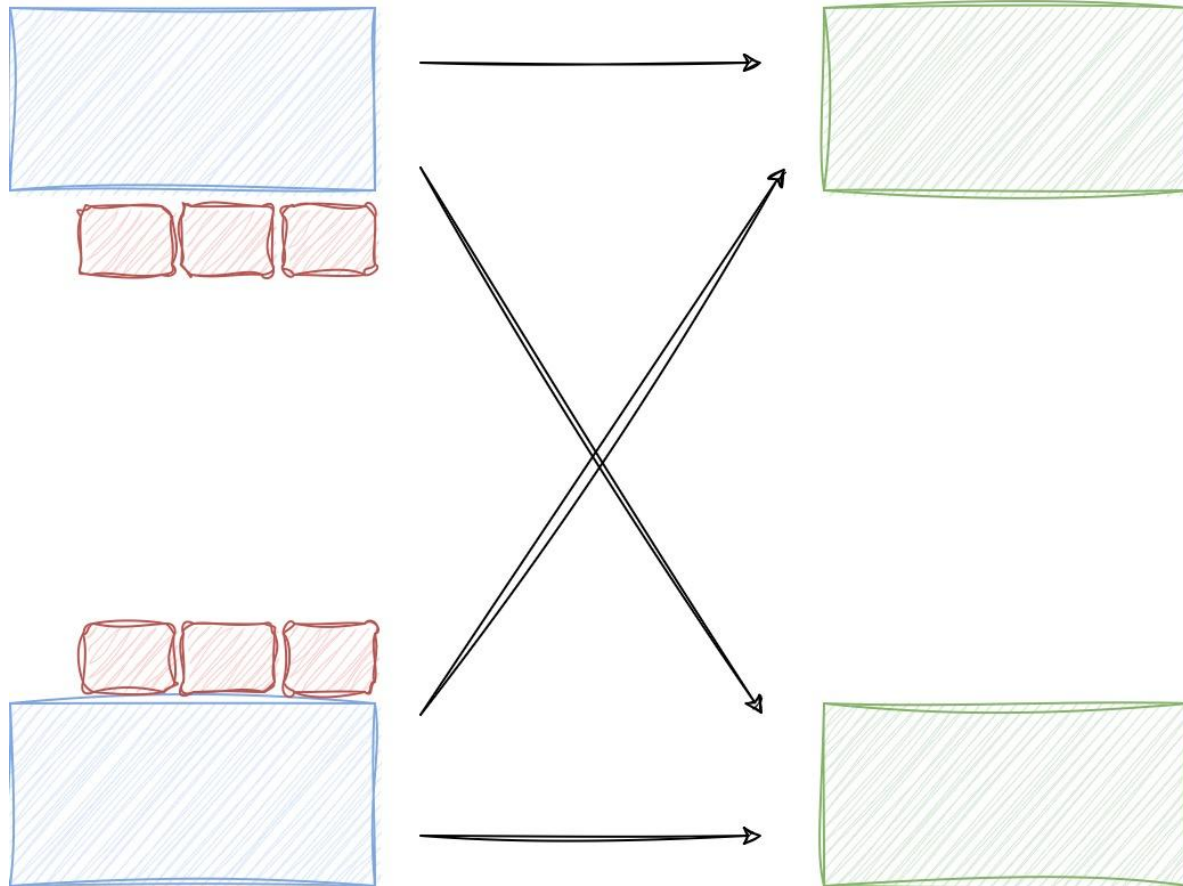
Scan

Aggregate



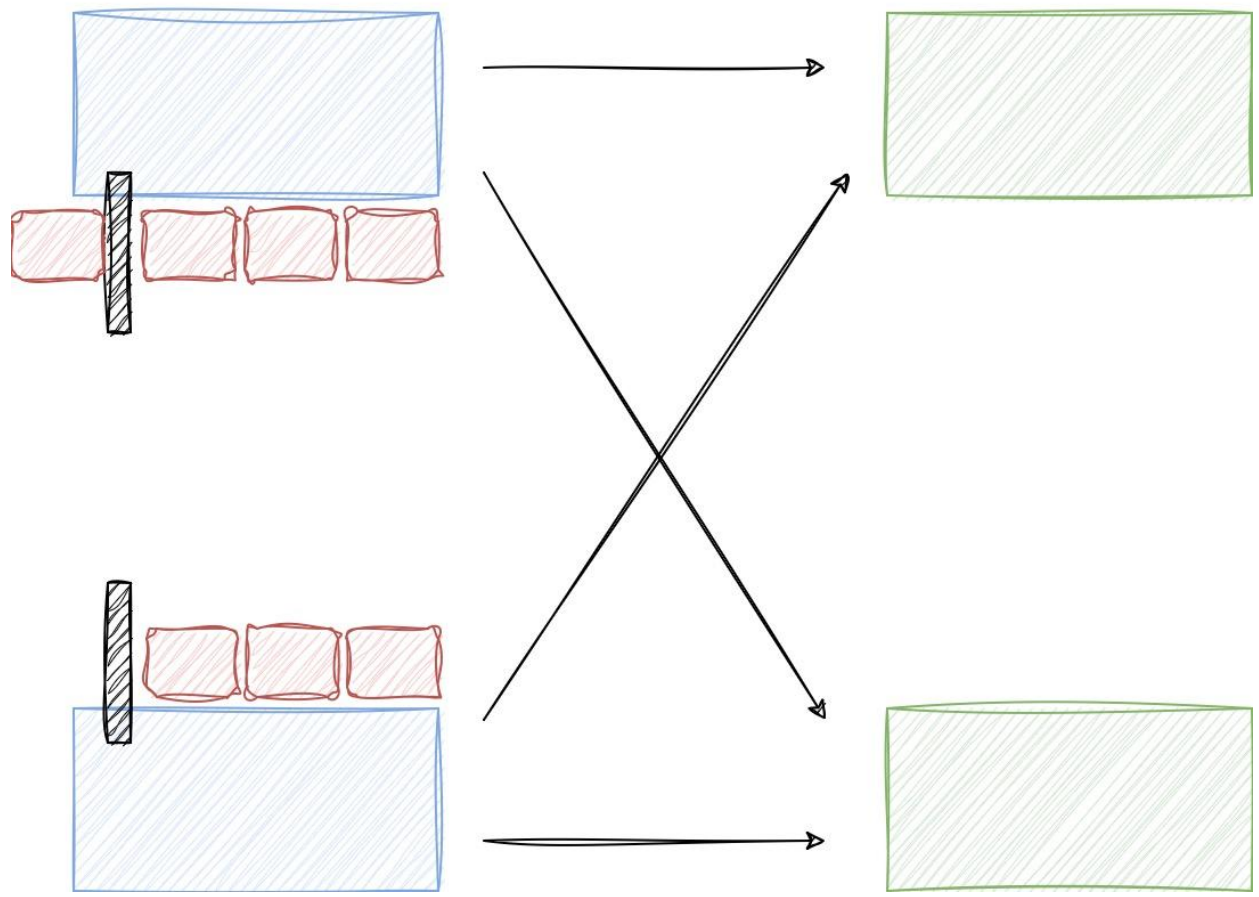
Scan

Aggregate



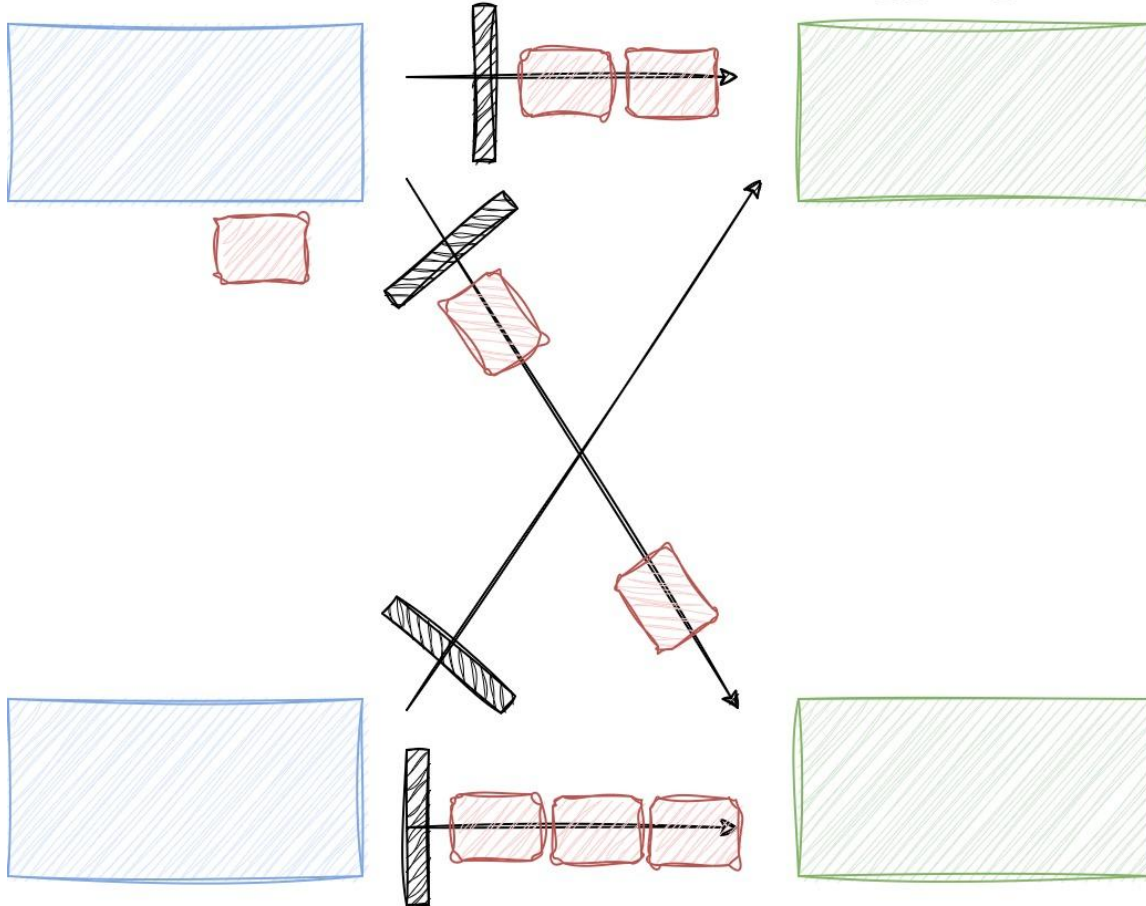
Scan

Aggregate

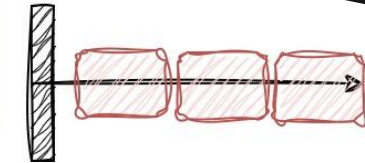


Scan

Aggregate



Scan



Aggregate

```
message BarrierMutation {  
  oneof mutation {  
    StopMutation stop = 4;  
    // Update outputs and hash mappings for some dispatchers, used for scaling.  
    UpdateMutation update = 5;  
    // Pause the dataflow of the whole streaming graph, only used for scaling.  
    PauseMutation pause = 7;  
    // Resume the dataflow of the whole streaming graph, only used for scaling.  
    ResumeMutation resume = 8;  
  }  
}
```

```
message Barrier {  
  enum BarrierKind {  
    BARRIER_KIND_UNSPECIFIED = 0;  
    // The first barrier after a fresh start or recovery.  
    BARRIER_KIND_INITIAL = 1;  
    // A normal barrier. Data should be flushed locally.  
    BARRIER_KIND_BARRIER = 2;  
    // A checkpoint barrier. Data should be synchronized to the shared storage.  
    BARRIER_KIND_CHECKPOINT = 3;  
  }  
  BarrierMutation mutation = 3;  
  data.Epoch epoch = 1;  
  //...  
}
```



Distributed systems Recovery

user	story_id
1234	1003
2345	1004
9999	1003
789	1004

story_id	COUNT(*)
1003	2
1004	2

story_id	COUNT(*)
1003	2
1004	1

ID	author
1003	Lu
1004	Jane
1005	John

Table:
Votes

Aggregate

- (1004, 1)
+ (1004, 2)

Join

MV

ID	author	title	url
1003	Lu	foo	http...
1004	Jane	bar	http...
1005	John	buzz	http...

Table:
Stories

ID	author	title	url	vcount
1003	Lu	foo	http	2
1004	Jane	bar	http	1

Event flow
direction



OP1



State A

OP2

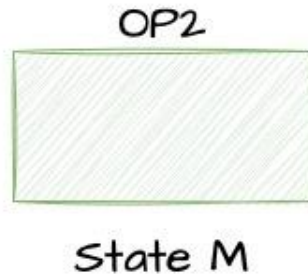
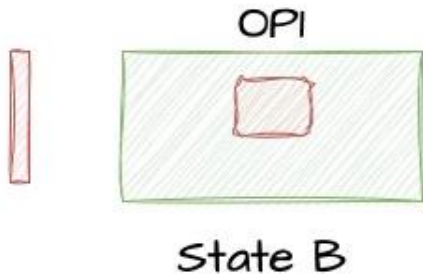


State M

Processing
Time: 0

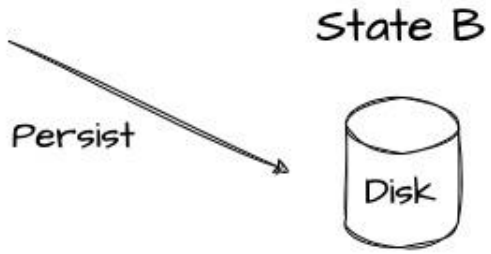
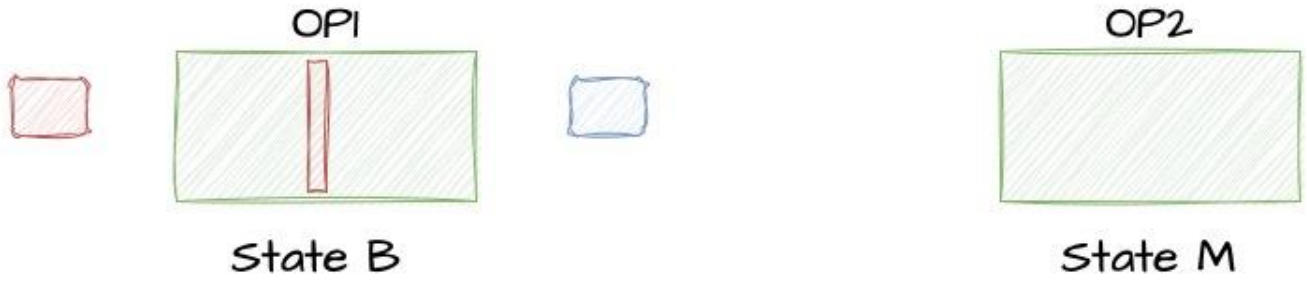
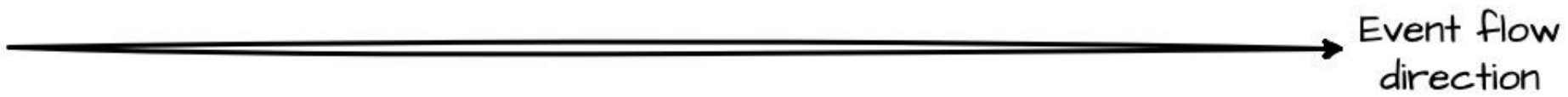


Event flow
direction



Processing
Time: I





Processing Time: 2

Event flow
direction



State C

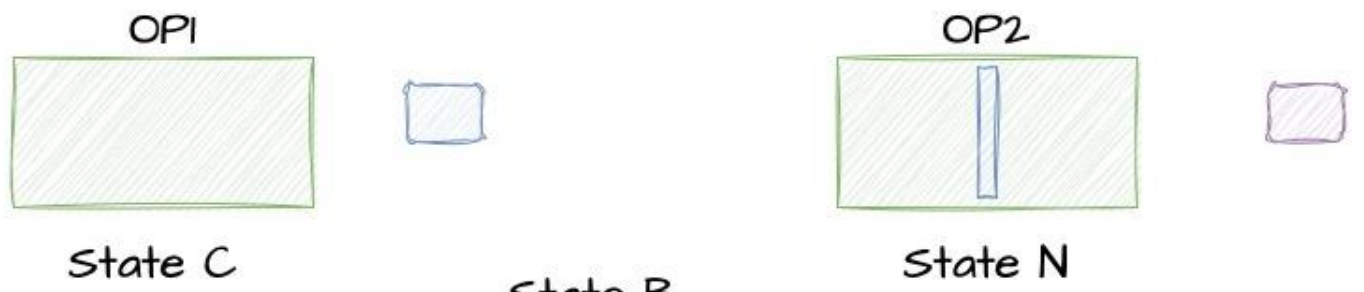
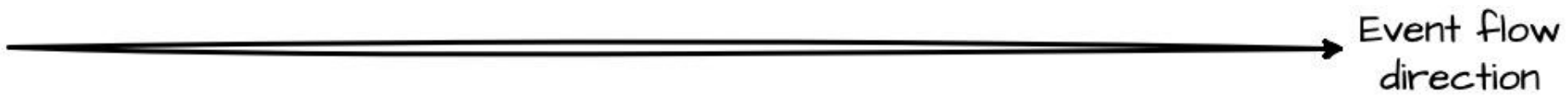
State N

State B

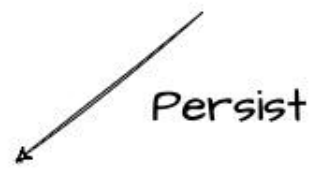


Disk

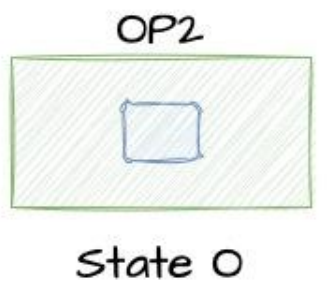
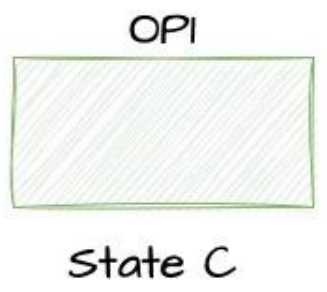
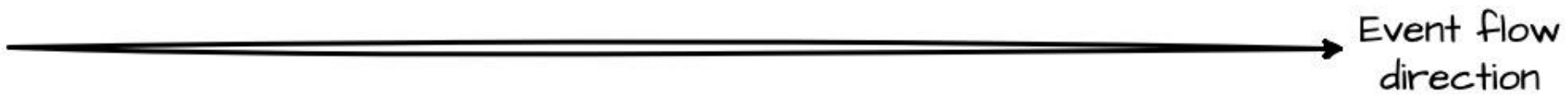
Processing
Time: 3



State B
State N



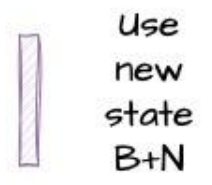
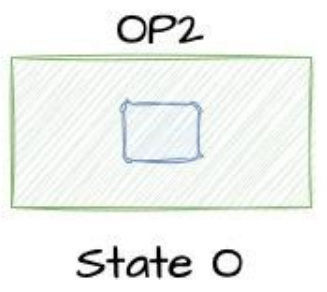
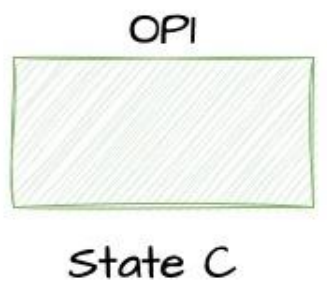
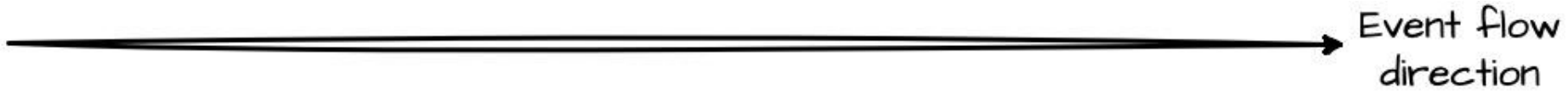
Processing
Time: 4



State B
State N



Processing
Time: 5



State B
State N

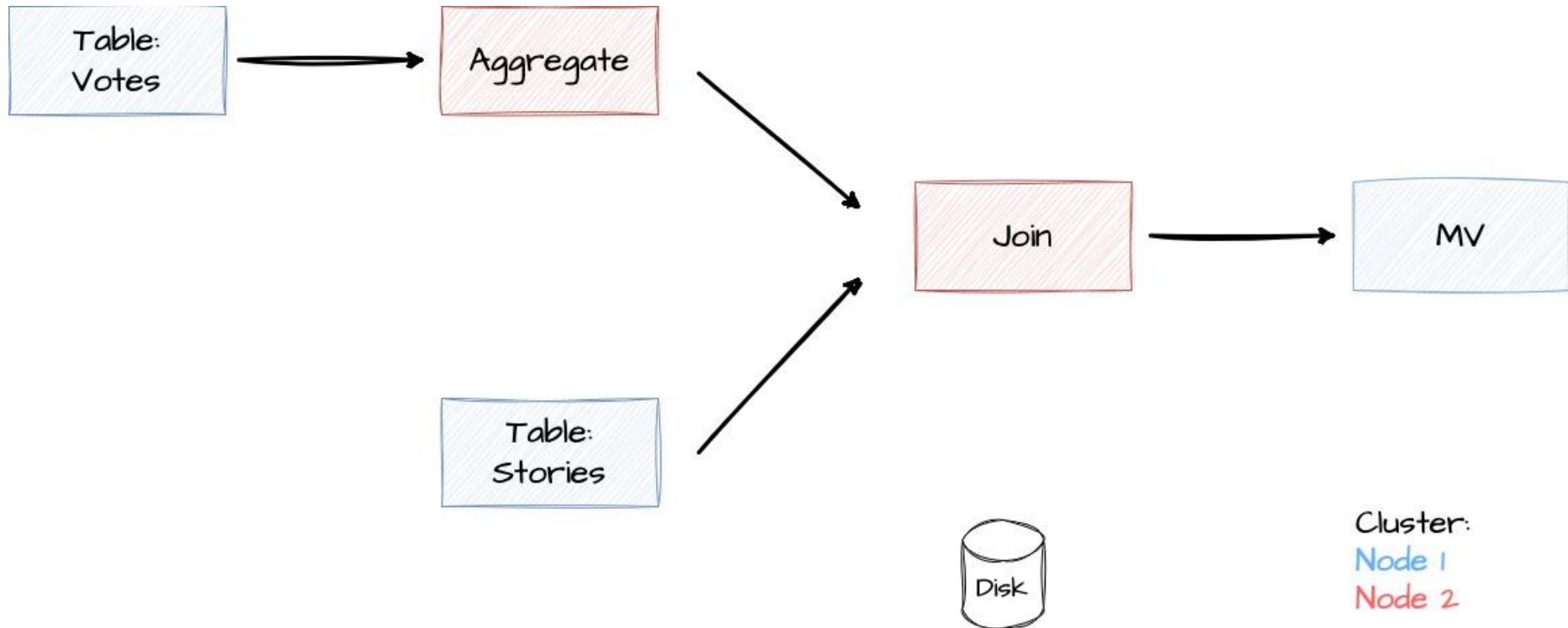


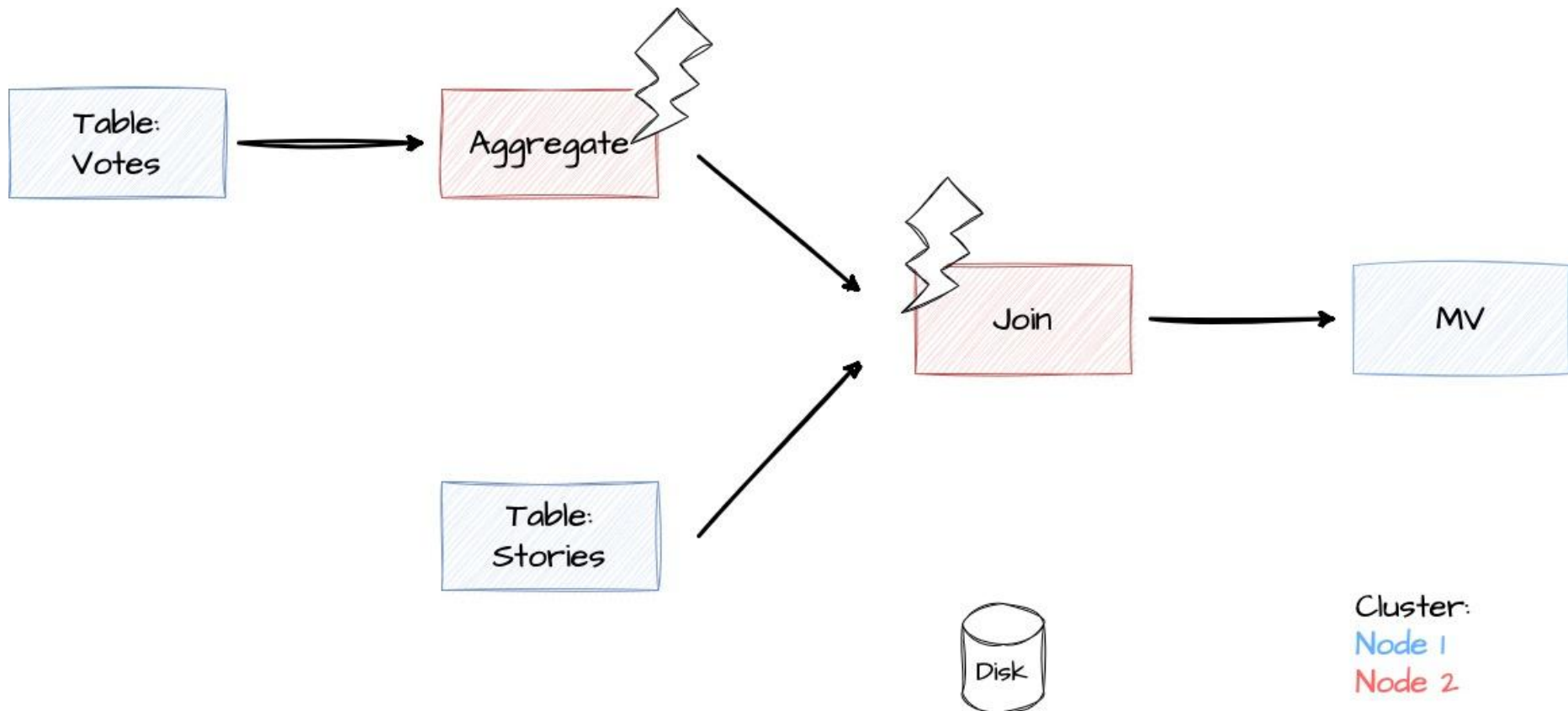
Processing
Time: 5

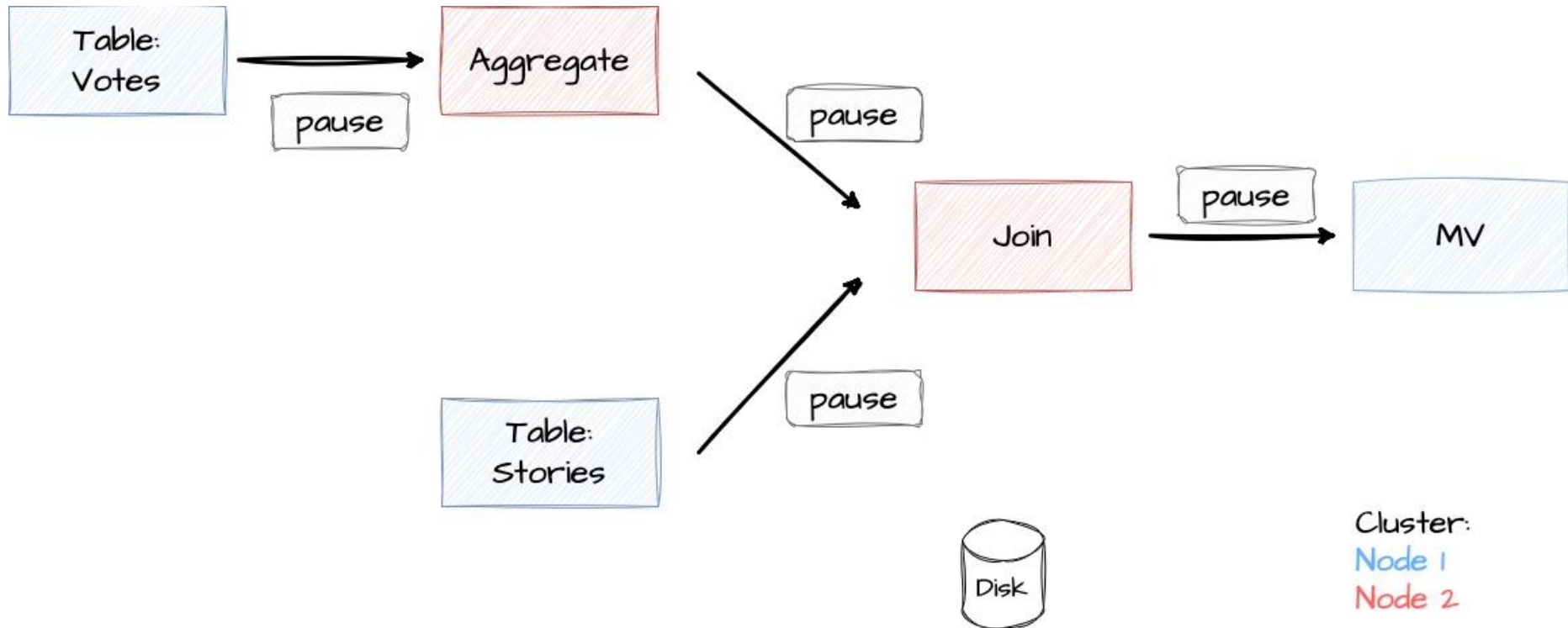


Distributed systems

Scaling







story_id	COUNT(*)
1003	2
1004	1

Table:
Votes

Aggregate

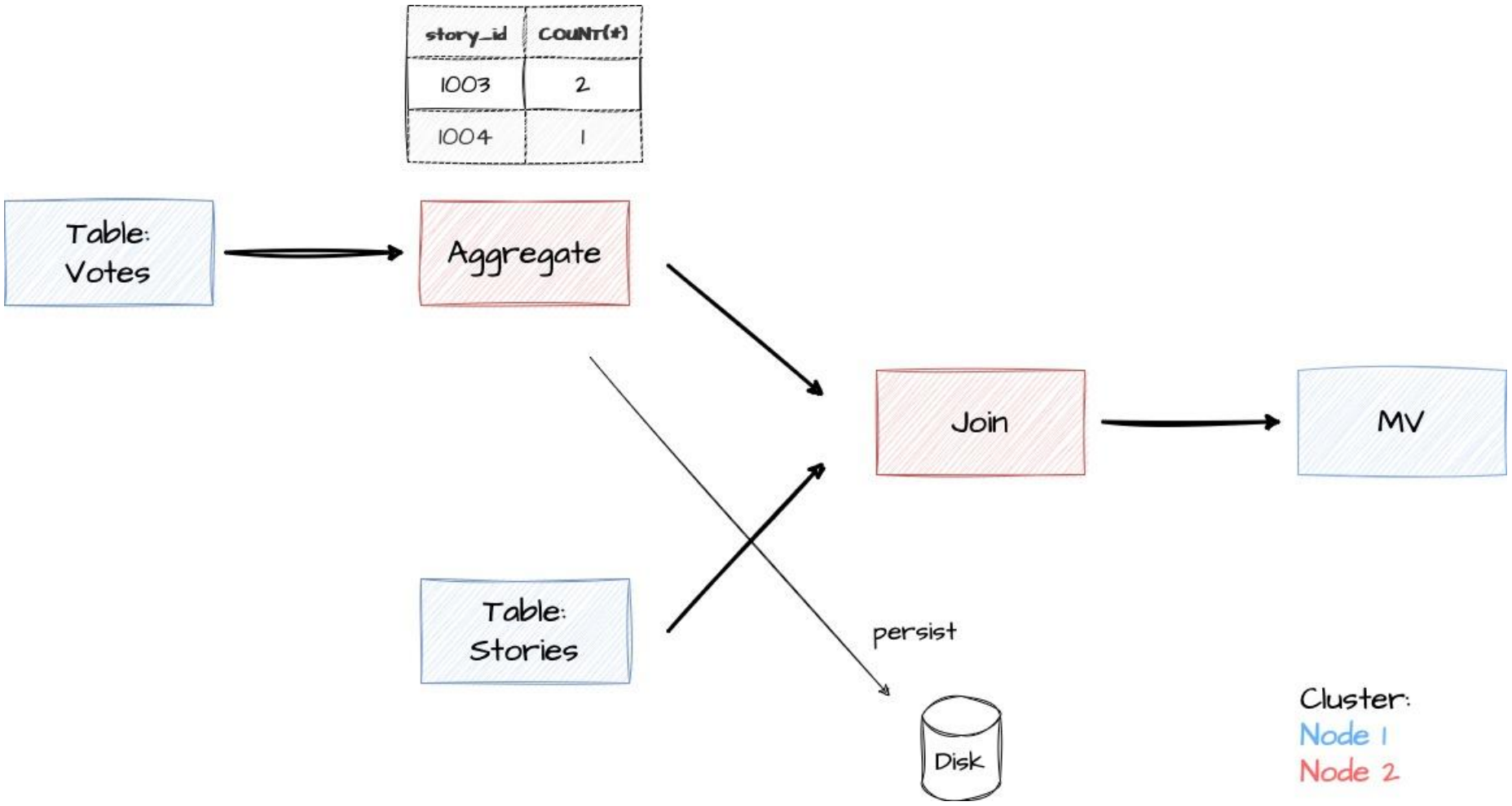
Table:
Stories

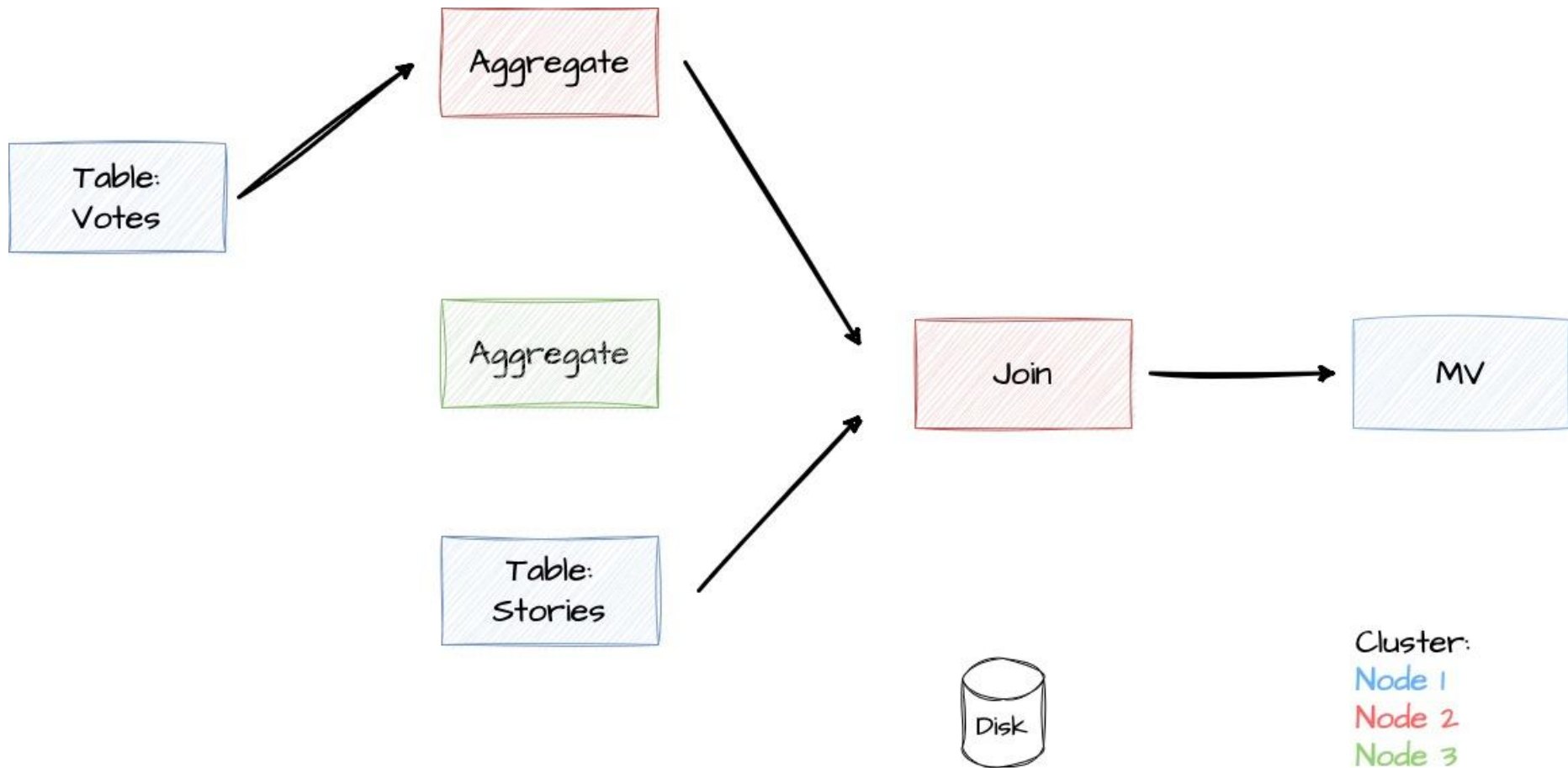
Join

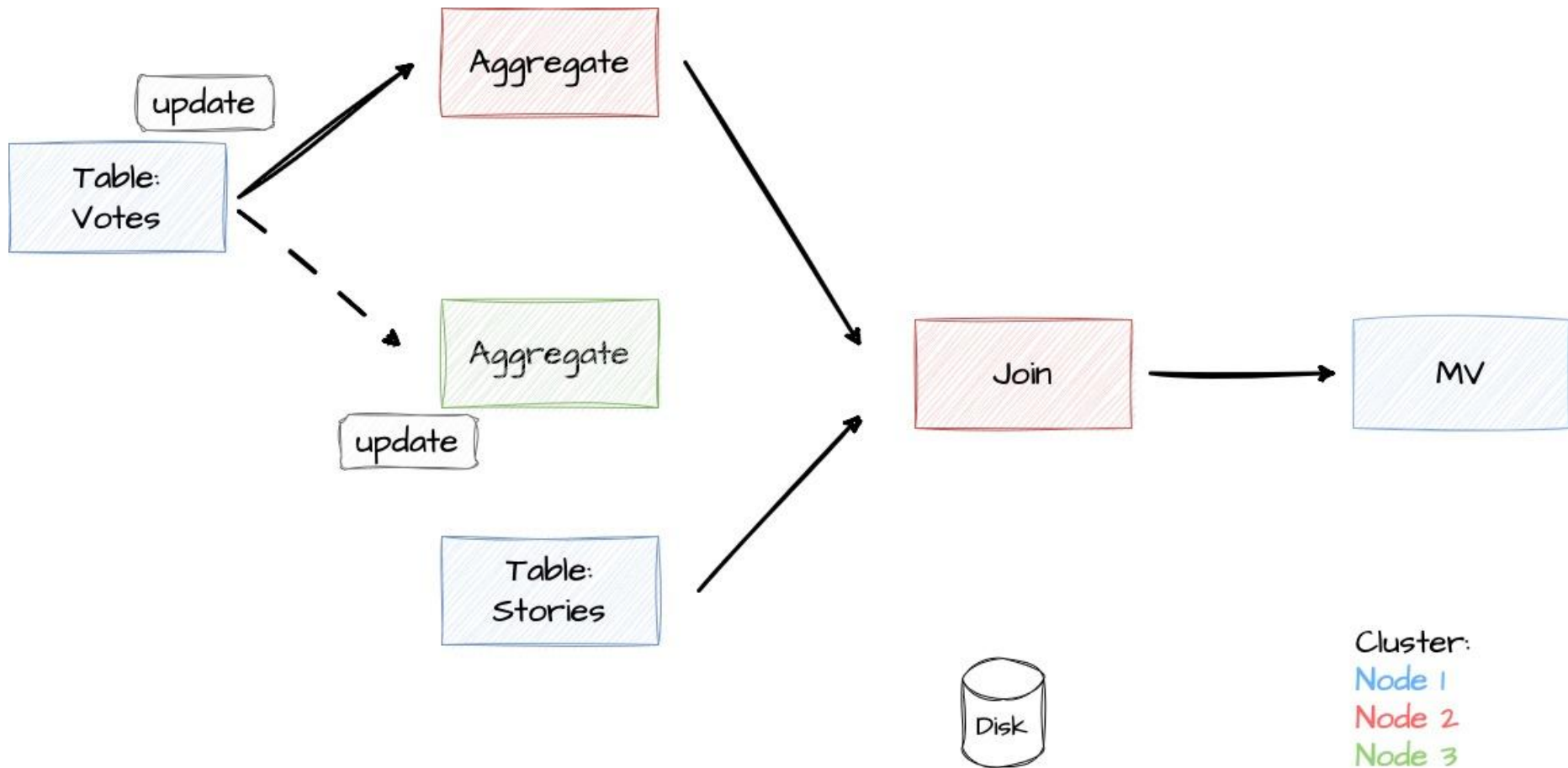
MV

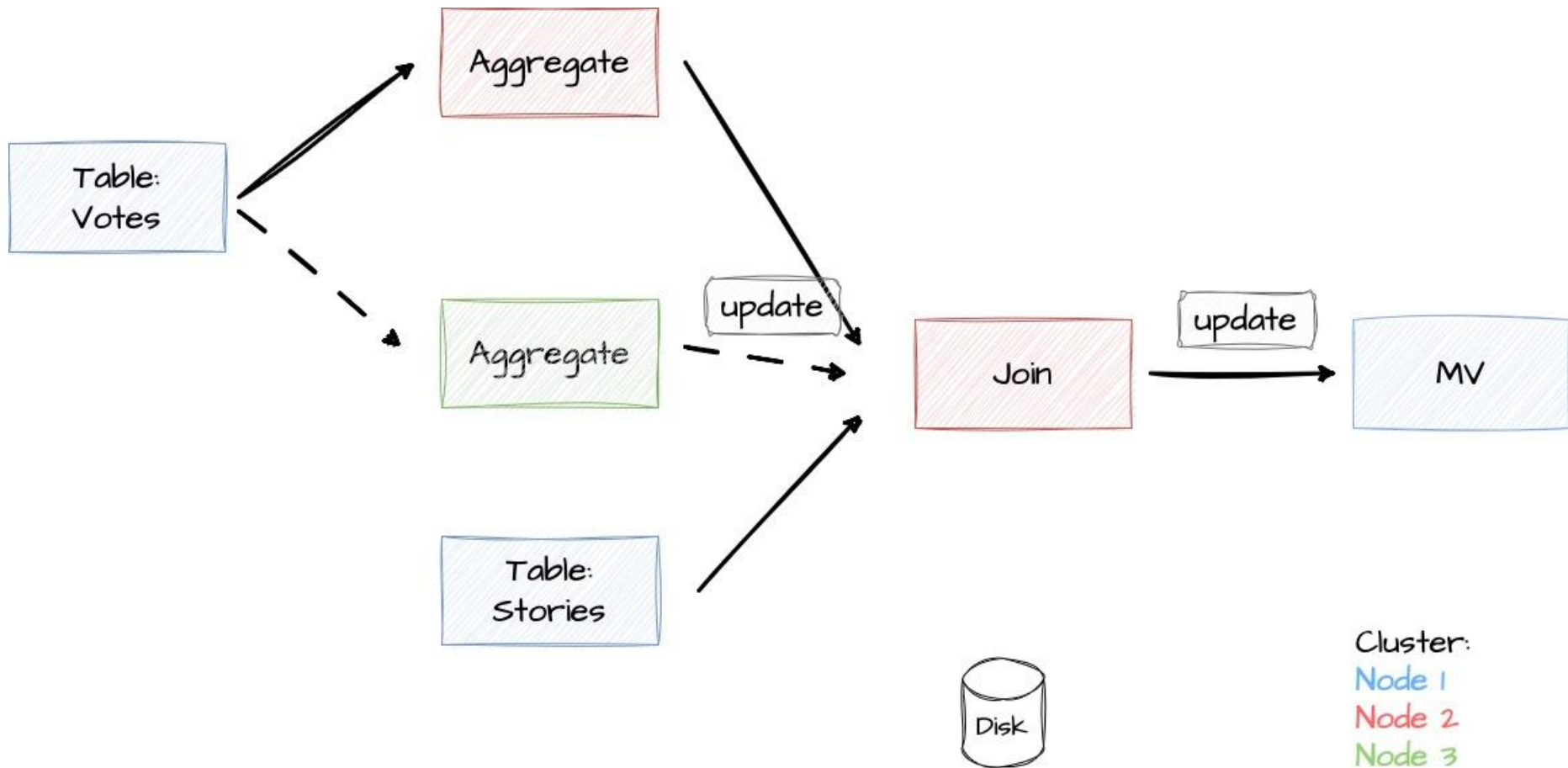


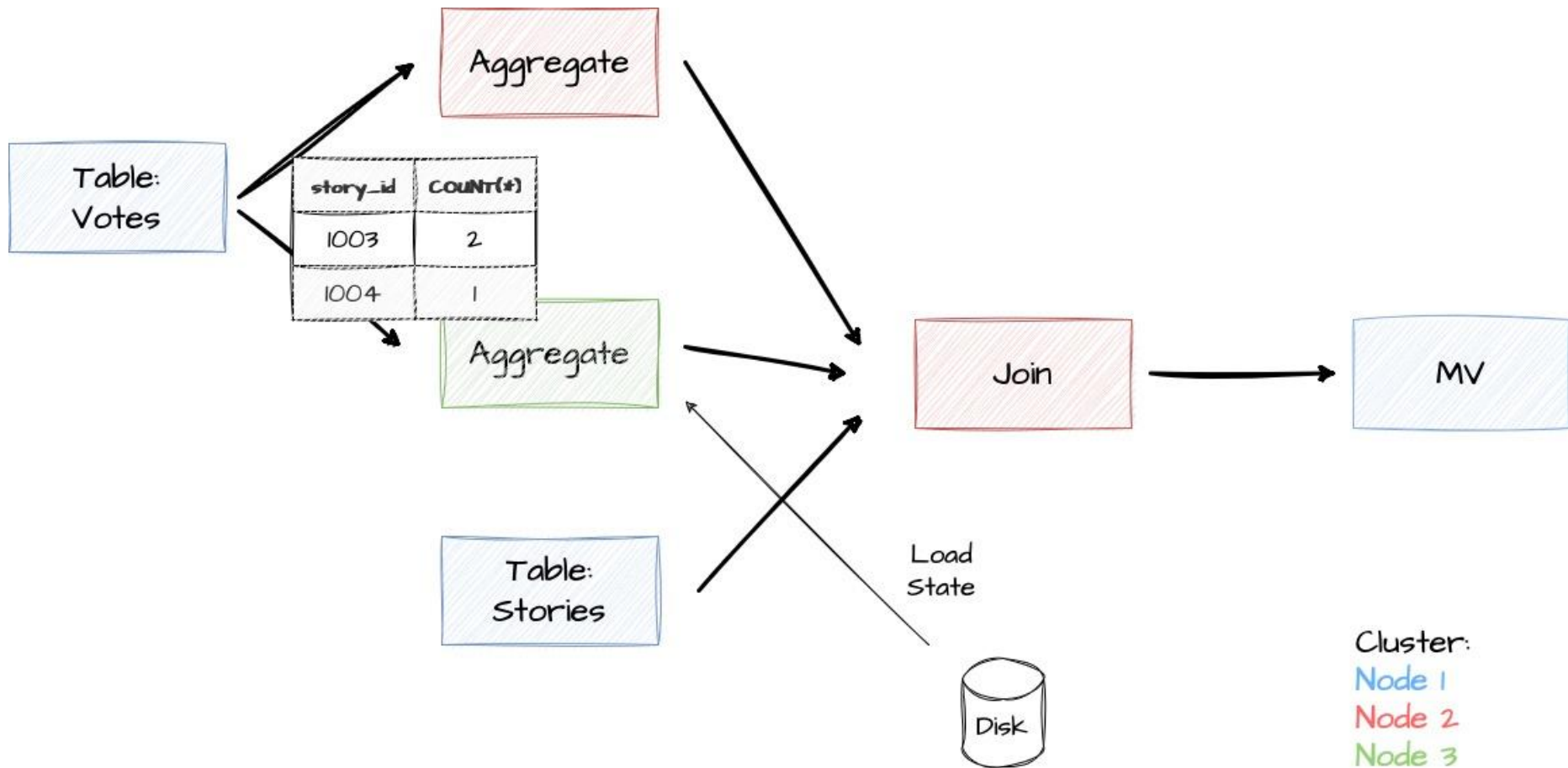
Cluster:
Node 1
Node 2

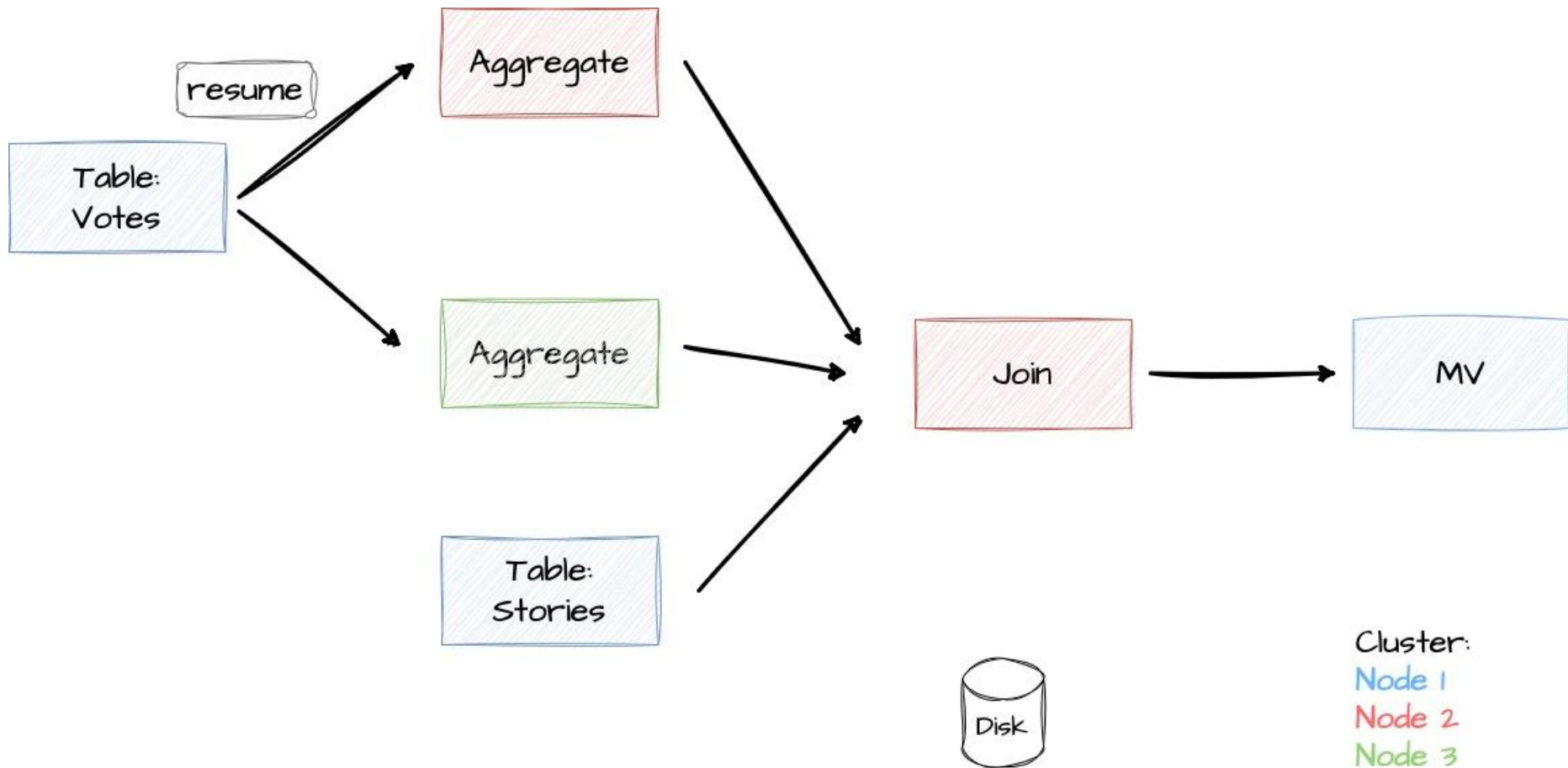














Thank you!



Try RisingWave:



Ask questions:

